ACOPhys - State University of St. Petersburg

Space-Filling Curves and Their Applications in Scientific Computing

Space-Filling Curves

Michael Bader

Technische Universität München, September 1-5, 2008





ТШП

Classification of Space-filling Curves

Definition: (*recursive* space-filling curve)

A space-filling curve $f: \mathcal{I} \to \mathcal{Q} \subset \mathbb{R}^n$ is called *recursive*, if both \mathcal{I} and \mathcal{Q} can be divided in *m* subintervals and sudomains, such that

- $f_*(\mathcal{I}^{(\mu)}) = \mathcal{Q}^{(\mu)}$ for all $\mu = 1, \dots, m$, and
- all $\mathcal{Q}^{(\mu)}$ are geometrically similar to \mathcal{Q} .

Definition: (contiguous space-filling curve)

A recursive space-filling curve is called *contiguous*, if for any two neighbouring intervals $\mathcal{I}^{(\nu)}$ and $\mathcal{I}^{(\mu)}$ also the corresponding subdomains $\mathcal{Q}^{(\nu)}$ and $\mathcal{Q}^{(\mu)}$ are direct neighbours, i.e. share an (n-1)-dimensional hyperplane.



Contiguous, Recursive Space-filling Curves

Examples:

- all Hilbert curves (2D, 3D, ...)
- also Peano curves and Sierpinksi curves

Properties: contiguous, recursive SFC are

- continuous (more exact: Hölder continuous with exponent 1/n)
- neighbourship-preserving
- describable by a grammar
- describable in an arithmetic form (see full set if slides)



3D Hilbert Curves

 Wanted: contiguous, recursive SFC, based on division-by-2 ⇒ leads to 3 basic patterns:



- in addition: symmetric forms, change of orientation
- always two different orientations of the components
- ⇒ numerous different Hilbert curves



3D Hilbert Curves – Iterations





3D Hilbert Curves – Variants

Different orientation of the sub-cubes:



- same basic pattern ("Motiv"), same approximating polygon
- differences only visible from 2nd iteration



Parallelisation using Space-filling Curves

Problem setting:

- "mesh" (2D, 3D, ...) of N unknowns ($N \gg 1000$)
- solve linear system(s) of equations (maybe repeatedly with varying right-hand side)
- in the system, only spatially neighbouring unknowns are coupled

Parallelisation:

Distribute N unknowns to p partitions, such that

- each partition contains the same number of unknowns (*load balancing*)
- for as many unknowns as possible, all neighbours are in the same partition (⇒ avoids communication bewtween partitions)



Parallelisation using Space-filling Curves (2)

Further demand: adaptivity

- add further unknowns (during/depending on intermediate results) or drop unknowns
- (re-)partitioning required to be **fast**: must not cost more computation time that going on with a bad load balance
- "shape preserving": if only few unknowns are added or dropped, the shape of partitions should not change strongly
 ⇒ only few unknowns then need to migrate to another partition
- \Rightarrow popular strategy: use **space-filling curves**



Hölder Continuity of Space-filling Curves

Definition: (Hölder continuous)

A function *f* is called *Hölder continuous with exponent r* on the interval *I*, if a constant C > 0 exists, such that for all $x, y \in I$:

 $||f(x) - f(y)||_2 \le C |x - y|^r$

Importance for space-filling curves:

- |x y| is the distance of the indices
- ||f(x) f(y)|| is the distance of the image points (in "space")
- To prove: the Hilbert curve is Hölder coninuous with exponent $r = d^{-1}$, where *d* is the dimension:

$$\|f(x) - f(y)\|_2 \le C |x - y|^{1/d} = C \sqrt[d]{|x - y|}$$

M. Bader: Space-Filling Curves and Their Applications in Scientific Computing ACOPhys – State University of St. Petersburg, September 1–5, 2008



пп

Hölder Continuity of the 3D Hilbert Curve

Proof analogous to simple continuity proof:

- given $x, y \in \mathcal{I}$; find an *n*, such that $8^{-(n+1)} < |x y| < 8^{-n}$
- 8^{-n} is the interval lenth for the *n*-tt iteration $\Rightarrow [x, y]$ covers at most two neighbouring(!) intervals.
- per construction of the 3D Hilbert curve, the function values h(x) and h(y) are in two adjacent cubes of side length 2⁻ⁿ.
- the length of the space diagonal through the two adjacent cubes is $2^{-n} \cdot \sqrt{1^2 + 1^2 + 2^2} = 2^{-n} \cdot \sqrt{6}$, hence:

$$\|h(x) - h(y)\|_{2} \le 2^{-n}\sqrt{6} = (8^{-n})^{1/3}\sqrt{6} = (8^{-(n+1)})^{1/3} 8^{1/3}\sqrt{6}$$

$$\le 2\sqrt{6} |x - y|^{1/3} \quad \text{q.e.d.}$$



Hölder Continuity and Parallelisation

• for the Hilbert curve (also Peano curve and all contiguous, recursive SFC), we have:

$$\|f(x)-f(y)\|_2 \leq C\sqrt[d]{|x-y|}$$

- relates the distance |x y| between indices to the disctance ||f(x) f(y)|| of (mesh) points
- |x y| corresponds to the section covered by the SFC (= area/volume)
- gives relation between volume (number of grid cells/points) and extent (e.g. radius) of a partition
- ⇒ Hölder continuity gives a quantitative estimate for compactness of partitions



Locality and Cartesian Grids

• from the Hölder continuity

$$\|f(x)-f(y)\|\leq C\sqrt[d]{|x-y|}$$

we directly obtain the following estimate:

$$\frac{\left\|f(x)-f(y)\right\|^{d}}{|x-y|} \leq \hat{C}$$

• for a discrete (cartesian, e.g.) grid, we can compare this to the relation

$$\frac{\left\|\mathbf{x}_{i}-\mathbf{x}_{j}\right\|^{d}}{\left|i-j\right|}\leq\hat{C},$$

where *i* and *j* are the indices (memory positions, e.g.) of two grid cells at positions \mathbf{x}_i and \mathbf{x}_j .

